

Overview of Linux I/O capabilities

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Certified Professional

Certified Professional

Oracle Database 12*c* Administrator Java SE 8 Programmer





Abakus Plus d.o.o.

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 - Services
 - OS & NET admin
 - DBA, Programming
 - Applications
 - Deja Vu
 - APPM
 - Arbiter

- Development Team
 - Enterprise Applications
 - Document Management
 - Newspaper Distribution
 - Flight Information System

References



Backup server

supports Oracle Databases and OLVM VMs

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Backup

takes no time

• Recovery

data recovery is almost instant

• Disk space

backed up data takes up minimal amount of disk space

• Availibility

data is always available and always in view

• Security

backed up data can not be deleted without support personnel intervention

• Alternative uses

BI analysis / reporting / DB upgrade verification / R&D testing / seamless business continuation

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APPM

Abakus Plus Performance Monitor

- For Oracle Database Standard Edition
- Made by DBAs for DBAs
- Temporal performance comparison
- Resource allocation optimization
- Database performance tracking
- Performance bottleneck optimization

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Linux ate my ram!

- https://www.linuxatemyram.com/
- Sometimes we do not want to populate the cache with one-time contents because we want other apps to keep their cached files.
- Great example would be a backup script

• Btw, from KVM hypervisor's perspective, populated page cache inside VM is used (RSS) memory.

Live Demo



live demo recap

Limiting page cache usage

- Cgroups
- sync; echo 3 > /proc/sys/vm/drop_caches
- dd if=random.iso iflag=<mark>nocache</mark> count=0

nocache uses POSIX_FADVISE to drop cache for whole file

POSIX_FADV_DONTNEED

\$ man fadvise, man posix_fadvise

#include <fcntl.h>

int posix_fadvise(int fd, off_t offset, off_t len, int advice);

Do not expect access in the near future. Subsequent access of pages in this range will succeed, but will result either in reloading of the memory contents from the underlying mapped file or zerofill-in-demand pages for mappings without an underlying file. live demo recap

POSIX_FADV_DONTNEED

	real	user	sys
always	<mark>9,12</mark>	0,65	<mark>8,19</mark>
after 16mb	8,87	0,60	8,05
after 32mb	8,88	<mark>0,61</mark>	8,00
after 64mb	8,63	0,62	7,80
after 256mb	8,37	0,58	7,63
never	7,37	0,62	6,36

FADV_SEQUENTIAL

POSIX_FADV_NORMAL sets the **readahead** window to the default size for the backing device; POSIX_FADV_SEQUENTIAL doubles this size, POSIX_FADV_RANDOM disables file readahead entirely.



stdio

- open()
- read()
- close()



Image from https://mattermost.com/blog/iouring-and-go/

uring



Images from https://mattermost.com/blog/iouring-and-go/

aio

- open()
- io_setup()
- io_prep_pread()
- io_submit()
- io_getevents()
- io_destroy()
- close()

uring

- open()
- io_uring_queue_init()
- io_uring_get_sqe()
- io_uring_prep_read()
- io_uring_sqe_set_data()
- io_uring_submit()
- io_uring_wait_cqe()
- io_uring_cqe_get_data()
- io_uring_cqe_seen()
- io_uring_queue_exit()
- close()

uring vs aio

- https://kernel.dk/io_uring.pdf
- aio only supports async IO for O_DIRECT
- aio may block for metadata operations
- uring is newer »replacement« for aio
 - more features, »faster«

- aio requires kernel ~2.6 or newer
- uring requires kernel ~5.6 or newer

Dirty Pages

- vm.dirty_background_[ratio|bytes] how many dirty pages before sync starts
- vm.dirty_[ratio|bytes] how many dirty pages before i/o is blocked until sync frees up the required space
- vm.dirty_expire_centisecs how long can a dirty page be in cache before sync starts
- vm.dirty_writeback_centisecs
 how often should kernel check if something needs
 to be done

fsync()

- fsync(), sync() causes all pending modifications to filesystem metadata and cached file data to be written to the underlying filesystems.
- fsync(int fd), syncfs(int fd) is like sync(), but synchronizes just the filesystem containing file referred to by the open file descriptor fd.

• Usage example: Oracle redolog files

sync_file_range()

SYNC_FILE_RANGE_WAIT_BEFORE | SYNC_FILE_RANGE_WRITE | SYNC_FILE_RANGE_WAIT_AFTER

will ensure that all pages in the specified range which were dirty when sync_file_range() was called are committed to disk.

fallocate() FALLOC_FL_PUNCH_HOLE

Specifying the FALLOC_FL_PUNCH_HOLE flag deallocates space (i.e., creates a hole) in the byte range starting at offset and continuing for len bytes.

Within the specified range, partial file system blocks are zeroed, and whole file system blocks are removed from the file. After a successful call, subsequent reads from this range will return zeroes.



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